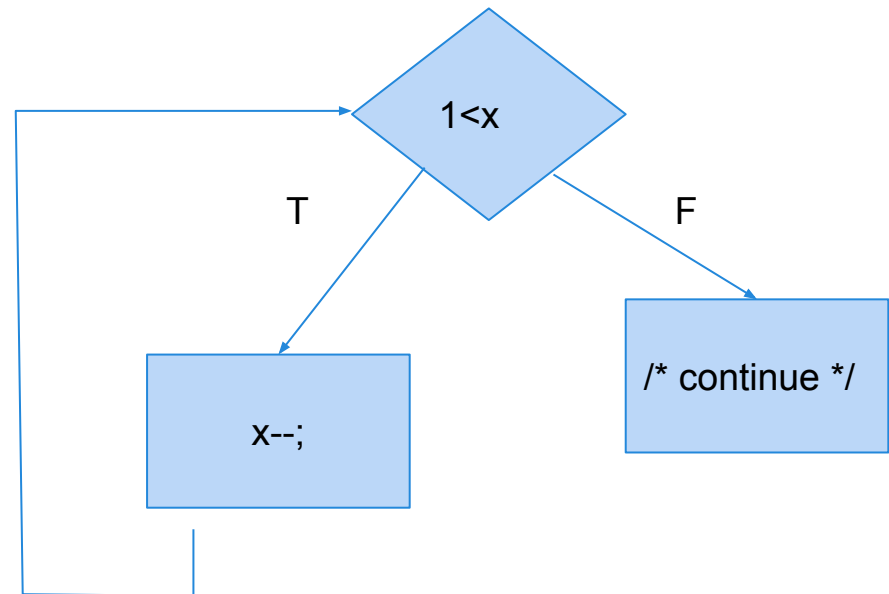


# Data Flow Testing

CSCE 747 - Lecture 10 - 02/09/2017

# Control Flow

- Capture dependencies in terms of how control passes between parts of a program.
- We care about the effect of a statement when it affects the path taken.
  - but deemphasize the information being transmitted.



# Data Flow

- Another view - program statements compute and transform data...
  - So, look at how that data is passed through the program.
- Reason about **data** dependence
  - A variable is used here - where does its value from?
  - Is this value ever used?
  - Is this variable properly initialized?
  - If the expression assigned to a variable is changed, what else would be affected?

# Data Flow Analyses

- Used to detect faults and other anomalies.

	<b>Any-Paths</b>	<b>All-Paths</b>
<b>Forward (pred)</b>	<b>Reach</b>  <i>U</i> may be preceded by <i>G</i> without an intervening <i>K</i>	<b>Avail</b>  <i>U</i> is always preceded by <i>G</i> without an intervening <i>K</i>
<b>Backward (succ)</b>	<b>Live</b>  <i>D</i> may lead to <i>G</i> before <i>K</i>	<b>Inevitability</b>  <i>D</i> always leads to <i>G</i> before <i>K</i>

- Also can be used to derive test cases.
  - Have we covered the data dependencies?

# Variable Aliasing

# Dealing With Arrays/Pointers

- Arrays and pointers (including object references and arguments) introduce issues.
  - It is not possible to determine whether two access refer to the same storage location.
    - `a[x] = 13;`  
`k = a[y];`
      - Are these a def-use pair?
    - `a[2] = 42;`  
`i = b[2];`
      - Are these a def-use pair?
        - Aliasing = two names refer to the same memory location.

# Aliasing

- *Aliasing* is when two names refer to the same memory location.

- `int[] a = new int[3];`

- `int[] b = a;`

- `a[2] = 42;`

- `i = b[2];`

- a and b are aliases.

- Worse in C:

- `p = &b;`

- `*(p + i) = k;`

# Uncertainty

- Dynamic references and aliasing introduce uncertainty into data flow analysis.
  - Instead of a definition or use of one variable, may have a potential def or use of a set of variables.
- Proper treatment depends on purpose of analysis:
  - If we examine variable initialization, might not want to treat assignment to a potential alias as initialization.
  - May wish to treat a use of a potential alias of  $v$  as a use of  $v$ .



# Dealing With Uncertainty

- Basic option: Treat all potential aliases as definitions and uses of the same variable:

```
a[1] = 13;           Def of a[1], use of a[2].  
k = a[2];
```

```
a[x] = 13;          Def and use of array a.  
k = a[y];
```

- Easiest and cheapest option when performing an analysis.
- Can be very imprecise.
  - They are only the same if  $x$  and  $y$  are the same.

# Dealing With Uncertainty

- Treat uncertainty about aliases like uncertainty about control flow.

```
a[x] = 13;  
k = a[y];
```

```
a[x] = 13;  
if(x == y)    k = a[x];  
else          k = a[y];
```

- In transformed code, all array references are distinct.
  - Any-path analysis - create a def-use pair, but assignment to `a[y]` does not erase definition to `a[x]`.
  - Gen sets include everything that might be references, kill sets only include definite references.

# Dealing With Uncertainty

```
a[x] = 13;  
k = a[y];
```

```
a[x] = 13;  
if(x == y)    k = a[x];  
else          k = a[y];
```

- In transformed code, all array references are distinct.
  - Any-path analysis - create a def-use pair, but assignment to `a[y]` does not erase definition to `a[x]`.
  - All-paths analysis - a definition to `a[x]` makes only that expression available. Assignment to `a[y]` kills `a[x]`.
    - Gen sets should include only what is definitely referenced and kill sets should include all possible aliases.

# Dealing With Nonlocal Information

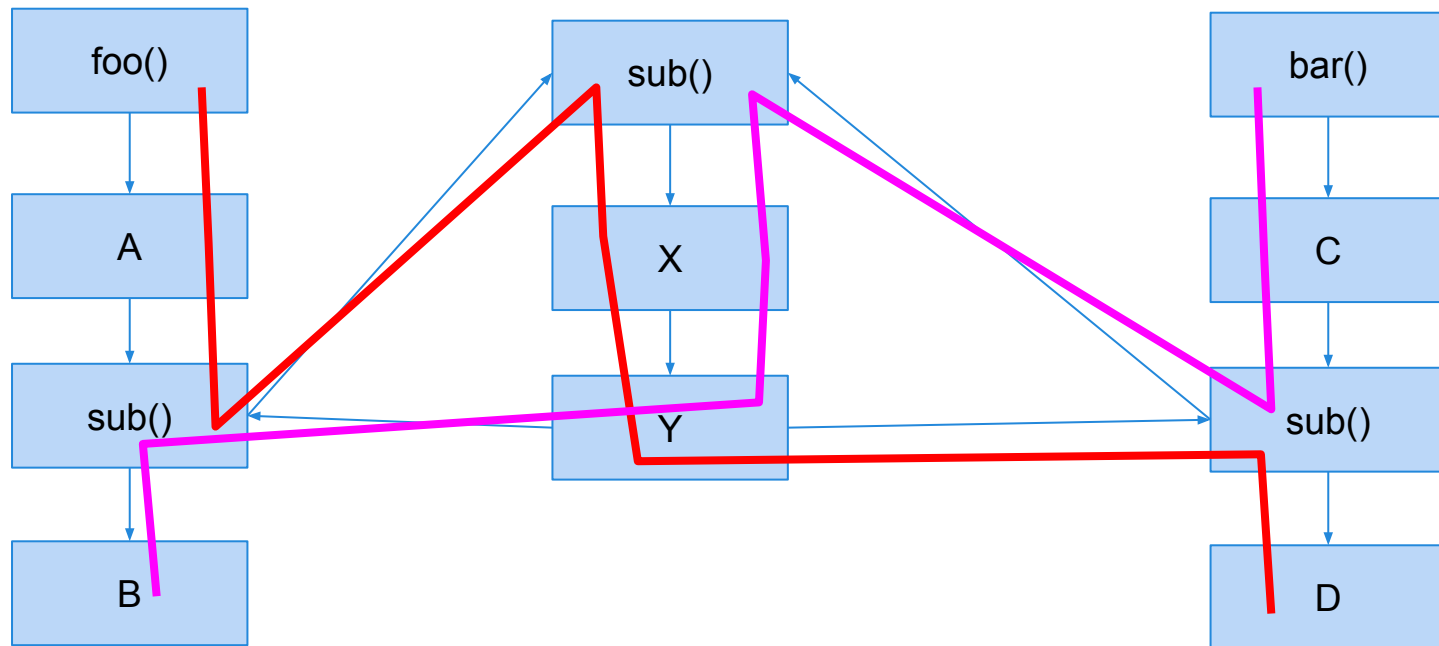
- fromCust and toCust may be references to the same object.
  - from/toHome and from/toWork may also reference the same object.
- Common option - treat all nonlocal information as unknown.
  - Treat Customer/PhoneNum objects as potential aliases.
  - Be careful - may result in results so imprecise they are useless.

```
public void transfer(Customer fromCust,
                    Customer toCust){
    PhoneNum fromHome =
        fromCust.getHomePhone();
    PhoneNum fromWork =
        fromCust.getWorkPhone();
    PhoneNum toHome =
        toCust.getHomePhone();
    PhoneNum toWork =
        toCust.getWorkPhone();
}
```

# Interprocedural Analysis

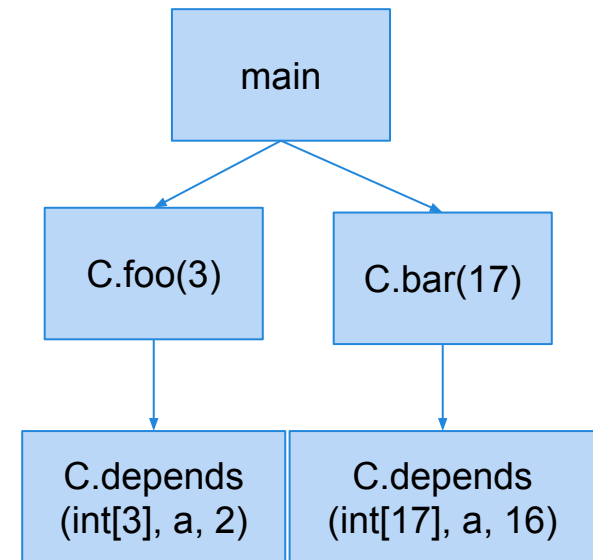
# Interprocedural Analysis - Control Flow

- First of **Problem - infeasible paths!** procedures in a large C



# Context-Sensitivity

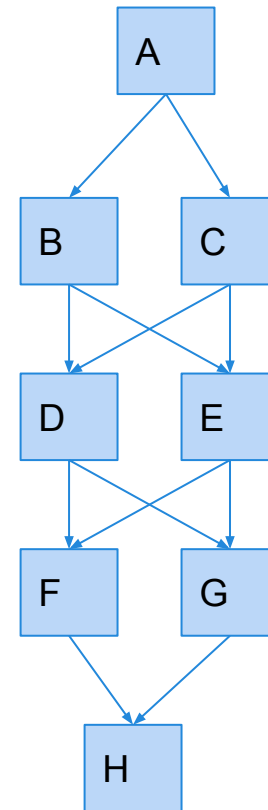
```
public class Context{
    public static void main(String args[]){
        Context c = new Context();
        c.foo(3);
        c.bar(17);
    }
    void foo(int n){
        int[] a = new int[n];
        depends(a,2);
    }
    void bar(int n){
        int[] a = new int[n];
        depends(a,16);
    }
    void depends(int[] a, int n){
        a[n] = 42;
    }
}
```



Context-Sensitive

# Context-Sensitive Analysis

- Copy the called procedure for each point that it is called.
- Problem - the number of contexts a procedure is called in is exponentially higher than the number of procedures.
  - Precise, but expensive analysis.
- In practice, only feasible for small groups of related procedures.



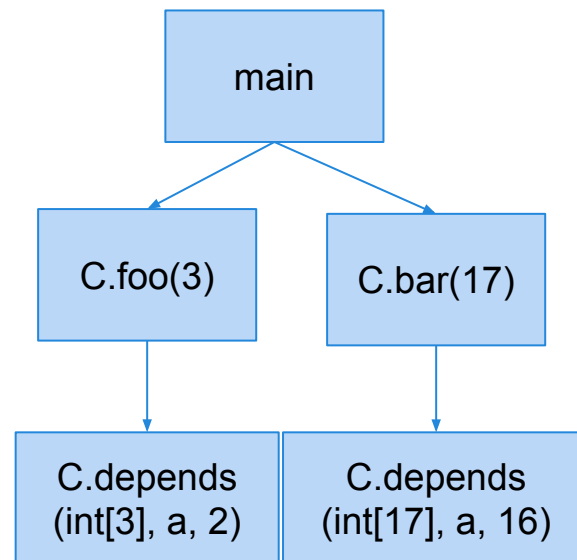


# Context-Insensitive Analysis

- Unhandled exception analysis
  - If procedure A calls procedure B that throws an exception, A must handle or declare that exception.
  - Analysis steps hierarchically through the call graph.
- Two conditions:
  - Information needed to analyze calling procedure must be small.
  - Information about the called procedure must be independent of caller (context-insensitive)
- Analysis can start from leaves of call graph and work upward to the root.

# Flow-Sensitivity

- Aliasing information requires context.
- Some analyses can sacrifice precision on another aspect: control-flow information
  - Call graphs are flow-insensitive.



# Insensitive Pointer Analysis

- Treat each statement as a constraint.  
 $x = y;$  (where  $y$  is a pointer)
- Note that  $x$  may refer to any of the same objects that  $y$  refers to.
  - $\text{References}(x) \supseteq \text{References}(y)$  is a constraint independent of the path taken.
  - Procedure calls are assignments of values to arguments.
- Results are imprecise, but better than just assuming that any two pointers might refer to the same object.

# Data Flow Testing

# Overcoming Limitations of Path Coverage

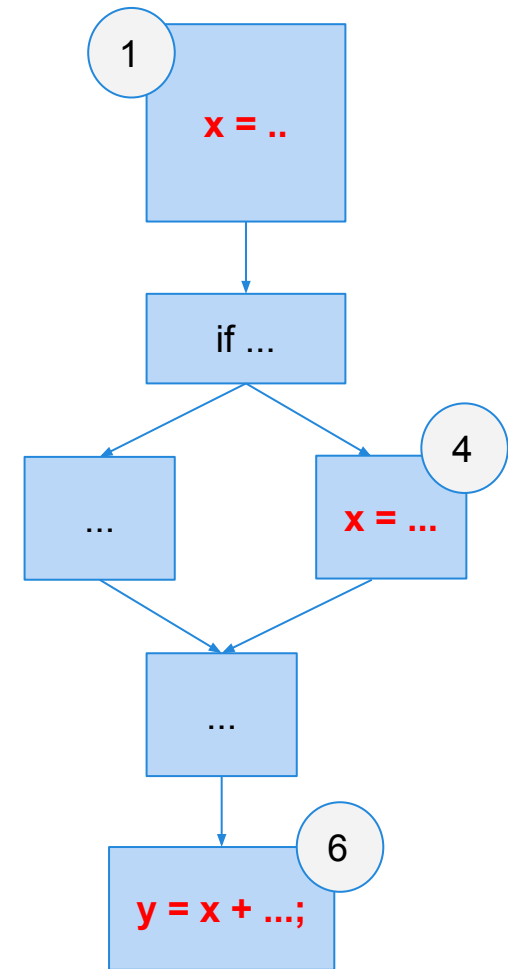
- We can potentially expose many faults by targeting particular *paths* of execution.
- Full path coverage is impossible.
- What are the important paths to cover?
  - Some methods impose heuristic limitations.
    - Loop boundary coverage
  - Can also use data flow information to select a subset of paths based on how one element can affect the computation of another.

# Choosing the Paths

- Branch or MC/DC coverage already cover many paths. What are the remaining paths that are important to cover?
- Basis of data flow testing - computing the wrong value leads to a failure only when that value is *used*.
  - Pair definitions with usages.
  - Ensure that definitions are actually used.
  - Select a path where a fault is more likely to propagate to an observable failure.

# Review - Def-Use Pairs

- Incorrect computation of  $x$  at either 1 or 4 could be revealed if used at 6.
- $(1,6)$  and  $(4,6)$  are *DU pairs* for  $x$ .
  - DU Pair = there exists a *definition-clear path* between the definition of  $x$  and a use of  $x$ .
  - If  $x$  is redefined on the path, the original definition is *killed* and replaced.



# Def-Use Pairs

- `++counter, counter++, counter+=1`  
`counter = counter + 1`
  - These are equivalent. They are a *use* of `counter`, then a new *definition* of `counter`.
- `char *ptr = *otherPtr`
  - Need a policy for how you deal with aliasing.
  - Ad-hoc option:
    - Definition of **string** `*ptr`
    - Use of **index** `ptr`, `string *otherPtr`, and `index otherPtr`.
- `ptr++`
  - Use of `index ptr`, and a definition of both the `index` and `string *ptr`.
  - Change to `index` moves the pointer to a new location.



# All DU Pair Coverage

- Requires each DU pair be exercised in at least one program execution.
  - Erroneous values produced by one statement might be revealed if used in another statement.

$$\text{Coverage} = \frac{\text{number exercised DU pairs}}{\text{number of DU pairs}}$$

- Can easily achieve structural coverage without covering all DU pairs.

# All DU Paths Coverage

- One DU pair might belong to many execution paths. Cover all simple (non-looping) paths at least once.
  - Can reveal faults where a path is exercised that should use a certain definition but doesn't.

$$\text{Coverage} = \frac{\text{number of exercised DU paths}}{\text{number of DU paths}}$$

# Path Explosion Problem

- Even without looping paths, the number of SU paths can be exponential to the size of the program.
- When code between definition and use is irrelevant to that variable, but contains many control paths.

```
void countBits(char ch){
    int count = 0;
    if (ch & 1)    ++count;
    if (ch & 2)    ++count;
    if (ch & 4)    ++count;
    if (ch & 8)    ++count;
    if (ch & 16)   ++count;
    if (ch & 32)   ++count;
    if (ch & 64)   ++count;
    if (ch & 128) ++count;
    printf(“‘%c’ (0X%02X) has %d bits
set to 1\n”, ch, ch, count);
}
```

# All Definitions Coverage

- All DU Pairs/All DU Paths are powerful and often practical, but may be too expensive in some situations.
- In those cases, pair each definition with at least one use.

$$\text{Coverage} = \frac{\text{number of covered definitions}}{\text{number of definitions}}$$

# Dealing With Aliasing

- Requires trade-off between precision and computational efficiency.
- Underestimate potential aliases
  - Could miss *def-use* pairs
- Overestimate potential aliases
  - Could have infeasible pairs, leading to unsatisfiable coverage obligations
- What is a suitable approximation of potential aliases for testing?

# Infeasibility Problem

- Metrics may ask for impossible test cases.
- Path-based metrics aggravates the problem by requiring infeasible combinations of feasible elements.
  - Alias analysis may add additional infeasible paths.
- All Definitions Coverage and All DU-Pairs Coverage often reasonable.
  - All DU-Paths is much harder to fulfill.

# Activity - DU Pairs

- Identify all DU pairs and write test cases to achieve All DU Pair Coverage.
  - Hint - remember that there is a loop.

```
1. int doSomething(int x, int y)
2. {
3.     while(y > 0) {
4.         if(x > 0) {
5.             y = y - x;
6.         }else {
7.             x = x + 1;
8.         }
9.     }
10.    return x + y;
11. }
```

# Activity - DU Pairs

```
1. int doSomething(int x, int y)
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9.     }
10.    return x + y;
11. }
```

Variable	Defs	Uses
x	1, 7	4, 5, 7, 10
y	1, 5	3, 5, 10

Variable	D-U Pairs
x	(1, 4), (1, 5), (1, 7), (1, 10), (7, 4), (7, 5), (7, 7), (7, 10)
y	(1, 3), (1, 5), (1, 10), (5, 3), (5, 5), (5, 10)



# Activity - DU Pairs

```
1. int doSomething(int x, int y)
2. {
3.     while(y > 0) {
4.         if(x > 0) {
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10.    return x + y;
11. }
```

Variable	D-U Pairs
x	<del>(1, 4), (1, 5), (1, 7), (1, 10),</del> <del>(7, 4), (7, 5), (7, 7), (7, 10)</del>
y	<del>(1, 3), (1, 5), (1, 10), (5, 3),</del> <del>(5, 5), (5, 10)</del>

**Test 1: (x = 1, y = 2)**

Covers lines 1, 3, 4, 5, 3, 4, 5, 3, 10

**Test 2: (x = -1, y = 1)**

Covers lines 1, 3, 4, 6, 7, 3, 4, 6, 7, 3, 4, 5, 3, 10

**Test 3: (x = 1, y = 0)**

Covers lines 1, 3, 8

# We Have Learned

- Arrays, pointers, and complex data structures introduce uncertainty into analysis.
  - Requires a policy for how aliasing is handled.
  - Trade-off between computational feasibility and precision.
- Analyses must handle non-local references.
  - Similar trade-off. Can gain efficiency by sacrificing flow sensitivity and context sensitivity.

# We Have Learned

- If there is a fault in a computation, we can observe it by looking at where the computation is used.
- By identifying DU pairs and paths, we can create tests that trigger faults along those paths.
  - All DU Pairs coverage
  - All DU Paths coverage
  - All Definitions coverage

# Next Class

- Model-Based Testing
- Reading: Chapter 14
- Homework:
  - Homework 2 is out - Due February 23
  - Reading Assignment 2 due tonight